Basics of Al and Machine Learning State-Space Search: Tree Search and Graph Search

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State-Space Search: Overview

Chapter overview: state-space search

- Foundations
- Basic Algorithms
 - Tree Search and Graph Search
 - Breadth-first Search
 - Uniform Cost Search
 - Depth-first Search
- Heuristic Algorithms

Introduction

Introduction

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General Search Algorithm

- Starting with initial state,
- repeatedly expand a state by generating its successors.
- Stop when a goal state is expanded
- or all reachable states have been considered.

Search Algorithms

General Search Algorithm

- Starting with initial state,
- repeatedly expand a state by generating its successors.
- Stop when a goal state is expanded
- or all reachable states have been considered.

In this chapter, we study two essential classes of search algorithms:

- tree search and
- graph search

(Each class consists of a large number of concrete algorithms.)

Tree Search

Tree Search

- possible paths to be explored organized in a tree (search tree)
- search nodes correspond 1:1 to paths from initial state
- duplicates (also: transpositions) possible, i.e., multiple nodes with identical state
- search tree can have unbounded depth

Generic Tree Search Algorithm

Generic Tree Search Algorithm

```
open := \mathbf{new} \ \mathsf{OpenList}
open.\mathsf{insert}(\mathsf{make\_root\_node}())
\mathbf{while} \ \mathbf{not} \ open.\mathsf{is\_empty}():
n := open.\mathsf{pop}()
\mathbf{if} \ \mathsf{is\_goal}(n.\mathsf{state}):
\mathbf{return} \ \mathsf{extract\_path}(n)
\mathbf{for} \ \mathbf{each} \ \langle a,s' \rangle \in \mathsf{succ}(n.\mathsf{state}):
n' := \mathsf{make\_node}(n,a,s')
open.\mathsf{insert}(n')
\mathbf{return} \ \mathsf{unsolvable}
```

discussion:

- generic template for tree search algorithms
- → for concrete algorithm, we must (at least) decide. how to implement the open list
 - concrete algorithms often conceptually follow template, (= generate the same search tree), but deviate from details for efficiency reasons

Graph Search

Reminder: Tree Search

reminder:

Tree Search

- possible paths to be explored organized in a tree (search tree)
- search nodes correspond 1:1 to paths from initial state
- duplicates (also: transpositions) possible, i.e., multiple nodes with identical state
- search tree can have unbounded depth

Graph Search

Graph Search

differences to tree search:

- recognize duplicates: when a state is reached on multiple paths, only keep one search node
- search nodes correspond 1:1 to reachable states
- search tree bounded, as number of states is finite

remarks:

- some graph search algorithms do not immediately eliminate all duplicates (→ later)
- one possible reason: find optimal solutions when a path to state s found later is cheaper than one found earlier

Generic Graph Search Algorithm

```
Generic Graph Search Algorithm
open := new OpenList
open.insert(make_root_node())
closed := new ClosedList
while not open.is_empty():
     n := open.pop()
     if closed.lookup(n.state) = none:
          closed.insert(n)
          if is_goal(n.state):
               return extract_path(n)
          for each \langle a, s' \rangle \in \text{succ}(n.\text{state}):
               n' := \mathsf{make\_node}(n, a, s')
               open.insert(n')
return unsolvable
```

Generic Graph Search Algorithm: Discussion

discussion:

- same comments as for generic tree search apply
- in "pure" algorithm, closed list does not actually need to store the search nodes
 - sufficient to implement closed as set of states
 - advanced algorithms often need access to the nodes, hence we show this more general version here
- some variants perform goal and duplicate tests elsewhere (earlier) → following chapters

Evaluating Search Algorithms

Criteria: Completeness

four criteria for evaluating search algorithms:

Completeness

Is the algorithm guaranteed to find a solution if one exists?

Does it terminate if no solution exists?

first property: semi-complete both properties: complete

Criteria: Optimality

four criteria for evaluating search algorithms:

Optimality

Are the solutions returned by the algorithm always optimal?

Criteria: Time Complexity

four criteria for evaluating search algorithms:

Time Complexity

How much time does the algorithm need until termination?

- usually worst case analysis
- usually measured in generated nodes

Criteria: Space Complexity

four criteria for evaluating search algorithms:

Space Complexity

How much memory does the algorithm use?

- usually worst case analysis
- usually measured in (concurrently) stored nodes

Summary

Summary

tree search:

search nodes correspond 1:1 to paths from initial state

graph search:

- search nodes correspond 1:1 to reachable states
- → duplicate elimination

generic methods with many possible variants